

Wyld Weyl and Twisted Bruhat: Partial Orders on Groups

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Introduction

Coxeter Systems (Coxeter Groups)

Classical Bruhat Order

Cocycles and "twisted Bruhat orders"

Current Research

Coxeter Groups

Setup: Let $S = \{s_1, s_2, \dots, s_n\}$ be a finite set. Suppose that we have a function $m : S \times S \rightarrow \mathbb{Z}_{\geq 1} \cup \infty$ satisfying

$$m(s_i, s_j) = m(s_j, s_i), \quad m(s_i, s_i) = 1 \quad \text{and} \quad m(s_i, s_j) \geq 2 \quad \text{for } i \neq j.$$

Definition

Let W be the group generated by S subject *only* to the relations $(s_i s_j)^{m(s_i, s_j)} = 1$ ($m(s_i, s_j) = \infty$ means there is no relation). In this case, we say that (W, S) is a **Coxeter system** (and we say W is a Coxeter group).

Coxeter Groups

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Implications

- W consists of “words” in S .
- For $s \in S$, $m(s, s) = 1$ means $ss = s^2 = 1$.
- For $r, s \in S$, $m(r, s) = k$ means

$$(rs)^k = \underbrace{rs \cdot rs \cdots rs \cdot rs}_{k \text{ times}} = 1$$

- This implies the “braid relations:”

$$\underbrace{rsr \cdots}_k = \underbrace{srs \cdots}_k$$

- For example, $m(r, s) = 2$ means $rs = sr$.

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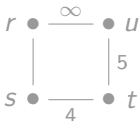
Organizing Data

Definition

A good way to organize this data is by using the so-called **Coxeter graph**. We construct the graph as follows:

1. Let the set of vertices be S .
2. We connect two vertices s_i and s_j if $m(s_i, s_j) \geq 3$.
3. We label the edge with $m(s_i, s_j)$ if $m(s_i, s_j) \geq 4$.

Example



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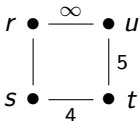
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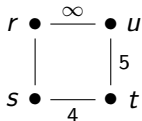
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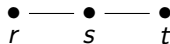
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S_3 : A Detailed Example

Standard Notation

1. Generators: (12) and (23)
2. $(12)(12) = 1$
 $(23)(23) = 1$
3. $(12)(23)(12) = (13)$
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Coxeter Group Notation

1. $r = (12)$ and $s = (23)$
2. $r^2 = 1$
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So, $S_3 = \langle r, s \mid r^2 = s^2 = (rs)^3 = 1 \rangle$

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Coxeter Graph:



Length

In S_3 , any word can be written with 3 letters or less:

$$rsrsrrrsr = rsrsr \underbrace{rr}_{1} sr = rsrs \underbrace{rsr}_{srs} = r \underbrace{srsrs}_{1} = r$$

Definition

For $w \in W$, if $w = s_{i_1} \dots s_{i_k}$ is written as short as possible, we say that this is a **reduced expression** for w and we say that w has **length** k . We write $\ell(w) = k$.

- $\ell : W \rightarrow \mathbb{N}$
- $\ell(1) = 0$
- w can have multiple reduced expressions
 - All have same length
 - e.g. $\ell(rsr) = \ell(srs)$ in S_3

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Posets

Definition

Given a set P , we say a binary relation \leq is a **partial order** on P if it satisfies three properties:

- Reflexivity: $a \leq a$ for all $a \in P$.
- Antisymmetry: If $a \leq b$ and $b \leq a$ then $a = b$.
- Transitivity: If $a \leq b$ and $b \leq c$ then $a \leq c$.

If we loosen the antisymmetry requirement, we call \leq a **pre-order**. A set with a partial order is called a **poset**.

Bruhat Order

We turn W into a poset in the following way:

- Fix an arbitrary reduced expression for $w \in W$. We say $v \leq w$ if v can be written as some subexpression of this reduced expression for w .
- 1 is the minimum element; poset is graded by length.
- For W finite, there is a unique maximum element.
- It is clearly a poset: reflexivity, antisymmetry, and transitivity are all clear via this definition.

Definition

The partial order defined above is called the **Bruhat order** on a Coxeter Group.

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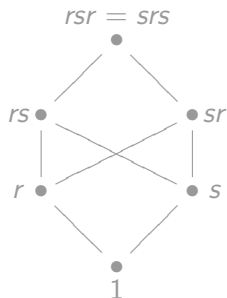
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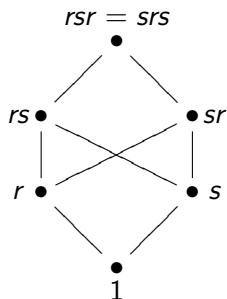
Bruhat Order for S_3



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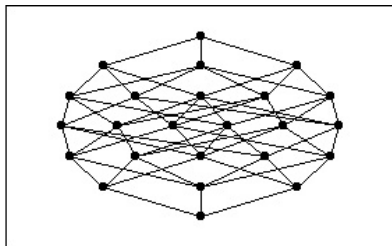
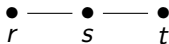
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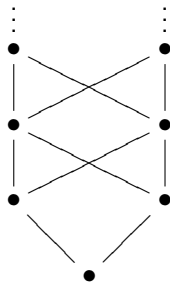
Some Pictures of Bruhat orders

Example (S_4)



<http://www.haverford.edu/math/cgreene/posets/posetgallery.html>

Example (D_∞)



A More Natural Definition (Cocycles)

Setup: Suppose (W, S) is a Coxeter System.

- $T = \bigcup_{w \in W} wSw^{-1}$, “reflections” of W .
- $\mathcal{P}(T)$ = power set of T
 - Forms a group under symmetric difference:

$$A + B = (A \cup B) \setminus (A \cap B)$$

- W acts on $\mathcal{P}(T)$ via conjugation:

$$wAw^{-1} \subseteq T \text{ for all } A \subseteq T$$

Definition

Using this action, a **cocycle** on W is a map $f : W \rightarrow \mathcal{P}(T)$ satisfying $f(xy) = f(x) + xf(y)x^{-1}$ for all $x, y \in W$.

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Using this action, a **cocycle** on W is a map $f : W \rightarrow \mathcal{P}(T)$ satisfying $f(xy) = f(x) + xf(y)x^{-1}$ for all $x, y \in W$.

A More Natural Definition (Bruhat Graph)

Definition

$N : W \rightarrow \mathcal{P}(T)$ given by $N(w) := \{t \in T \mid \ell(tw) < \ell(w)\}$ is called the **reflection cocycle**.

We can “twist” the conjugation action with this cocycle to get a new action: For $A \in \mathcal{P}(T)$ and $w \in W$

$$w \cdot A = N(w) + wAw^{-1}$$

Definition

Let Ω_W to be the directed graph given by the following:

1. The vertices are the elements of W .
2. The edge set is $E_W = \{(tw, w) \mid t \in T \text{ and } t \in w \cdot \emptyset\}$.

Then, the Bruhat order can be restated as $v \leq w$ iff there is a path from v to w in Ω_W .

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Generalization

Definition

For any $A \in \mathcal{P}(T)$, let $\Omega_{(W,A)}$ be the directed graph given by the following:

1. The vertices are the elements of W .
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Question

When is \leq_A a partial order? In other words, for what $A \in \mathcal{P}(T)$ is \leq_A antisymmetric?

Subsets of T

Definition

A **dihedral reflection subgroup** is a subgroup of W generated by 2 elements of T . Any dihedral reflection subgroup has a canonical set of generators.

- D is a dihedral reflection subgroup generated by r and s .

$$T \cap D = \{r, rsr, rsrsr, \dots, srsrs, srs, s\}$$

Definition

We say total order \prec on T is a **reflection order** if one of

- $r \prec rsr \prec \dots \prec srs \prec s$
- $s \prec srs \prec \dots \prec rsr \prec r$

is true for any dihedral reflection subgroup D of W .

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The Main Result

Definition

A subset $A \subseteq T$ is an **initial section of a reflection order** (i.s.o.r.o) if there is some reflection order \prec with the property that $a \prec b$ for all $a \in A$ and $b \in T \setminus A$.

Now, we define

$$\mathbf{A}_{(W,S)} = \{A \subseteq T \mid A \cap D \text{ is an i.s.o.r.o on } D\}$$

where D ranges over all dihedral reflection subgroups of W .

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Theorem

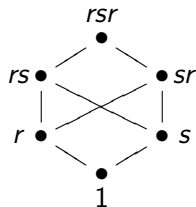
Let (W, S) be a Coxeter system with reflections T and let $A \subseteq T$. \leq_A is a partial order if and only if $A \in \mathbf{A}_{(W,S)}$.

S_3 : A Detailed Example

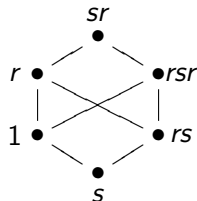
- S_3 is a dihedral group (and has no proper dihedral subgroups).
- $T = \{r, rsr, s\}$; $\mathcal{P}(T)$ has 8 elements.
- 2 possible reflection orders: $r \prec rsr \prec s$ or $s \prec rsr \prec r$.
- 6 initial sections: $\emptyset, \{r\}, \{s\}, \{r, rsr\}, \{s, rsr\}, T$.
- 2 non-initial sections: $\{r, s\}$ and $\{rsr\}$.
- 6 twisted Bruhat orders. (All isomorphic since S_3 is finite.)

S_3 : A Detailed Example

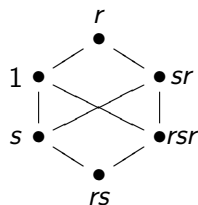
$$A = \emptyset$$



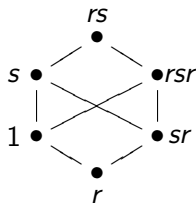
$$A = \{s\}$$



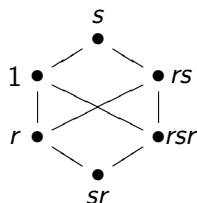
$$A = \{s, rsr\}$$



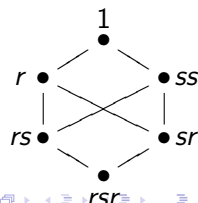
$$A = \{r\}$$



$$A = \{r, rsr\}$$



$$A = T = \{r, rsr, s\}$$



Current Research

- Representation Theory associated to twisted Bruhat orders:
 - Kazhdan-Lusztig Polynomials.
 - Positivity conjectures.
- Classify $H^1(W, \mathcal{P}(T))$ for Coxeter groups (Finite, Weyl, etc.).
- Do other cocycles lead to posets on W in a similar manner?
- Use cocycles to understand other properties of Coxeter groups.

Thanks!

Any Questions?



Bruhat



Weyl